Heterogeneous Edge Caching based on Actor-Critic Learning with Attention Mechanism Aiding

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Abstract—In recent years, the explosive growth of network traffic has placed significant strain on backbone networks. To alleviate the content access delay and reduce additional network resource consumption resulting from large-scale requests, edge caching has emerged as a promising technology. Despite capturing substantial attention from both academia and industry, most existing studies overlook the heterogeneity of the environment and the spatial-temporal characteristics of content popularity. As a result, the potential for edge caching remains largely unexploited. To address these challenges, we propose a neighborhood-aware caching (NAC) framework in this paper. The framework leverages the perimeter information from neighboring base stations (BSs) to model the edge caching problem in heterogeneous scenarios as a Markov Decision Process (MDP). To fully exploit the environmental information, we introduce an improved actor-critic method that integrates an attention mechanism into the neural network. The actor-network in our framework is responsible for making caching decisions based on local information, while the critic network evaluates and enhances the actor’s performance. The multi-head attention layer in the critic network enables integration of environmental features into the model, reducing the limitations associated with local investigation. To facilitate comparison from an engineering perspective, we also propose a heuristic algorithm, Neighbor-Influence-Least-Frequently-Use (NILFU). Our extensive experiments demonstrate that the proposed NAC framework outperforms other baseline methods in terms of average delay, hit rate, and traffic offload ratio in heterogeneous scenarios. This highlights the effectiveness of the neighborhood-aware caching approach in enhancing the performance of edge caching systems in such scenarios.

Index Terms—Edge Caching, Attention Mechanism, Actor-Critic Learning, Multi-agent Caching

I. INTRODUCTION

Recently, the development momentum of network technology is swift and skyrocketing, which leads to the deepening of digital transformation in all walks of life. According to Cisco’s forecast, global data traffic will increase three times faster than it did four years ago [1]. The enormous data traffic has become a severe challenge for mobile network operators (MNOs), pushing the urgent revolutions of network architecture and high-level communication technologies (e.g., 5G). As a promising technology, mobile edge caching (MEC) is widely used in edge networks (HetENets), which are densely designed to the edge and high-level communication technologies (e.g., 5G). As a promising technology, mobile edge caching (MEC) is widely conducted in ICTFICIAL OY and is partially supported by the European Union’s Horizon Europe program for Research and Innovation through the aerOS project under Grant No. 101069732. It was also partially supported by the Academy of Finland 6Genesis project under Grant No. 318927 and the Academy of Finland IDEA-MILL project under Grant No. 352428.

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enlighten the environmental awareness capacity of the network system [4]. Based on the previous discussions, there are several challenges that need to be considered as follows.  

- **Dynamic Environment is Changing**: Existing studies are hard to make adjustment decisions during the information exchange of environmental interactions.  
- **Contextual Information is Ignored**: Contextual information from neighbor BSs is always ignored, leading to low network performance when designing the caching strategy in the heterogeneous environment.  

The attention mechanism is widely used to fix this gap since it is proficient in blending the elements according to their degree of influence [6]. By using the attention mechanism, the latent influences between the entities can be found. This helps reach the goal of coordination and balance from different points of view at the same time. Inspired by the MAAC algorithm proposed in the work [7], which selects appropriate information to estimate the execution effect of actors based on a multi-agent system by attention mechanism, we propose a multi-agent neighbor-aware actor-critic framework, nominated NAC. In this method, we consider not only the knowledge of the current area but also the global information of neighboring nodes. The actor-network in our framework is responsible for making caching decisions based on the local information, while the critic network is responsible for evaluating and correcting the actor’s performance. The introduced multi-head attention layer in the critic network helps integrate the environment features into the model and eliminate the limitations caused by the local perspective. The main contributions of this paper are listed as follows:

- To address the problems of a dynamic environment and limited cache resources, we utilize a multi-head attention technique to effectively capture system changes and the status of neighbor BSs, therefore addressing the issue of adaptation and cooperation.
- We model the heterogeneous edge caching issue as the Markov Decision Process (MDP) and introduce the NAC framework for neighbor-aware actor-critic collaborative edge caching. In addition, we develop an engineering solution based on information by enhancing the classic Least-Frequently Used (LFU) algorithm for scenarios that cannot be integrated by neural network capabilities.
- Simulation results indicate that, compared to current baseline approaches, the proposed framework can effectively enhance hit rate, and decrease average transmission latency, and traffic offload ratio in a heterogeneous environment.

The rest of this paper is organized by the following: Section II introduces the related studies in the aspects of edge caching optimization and attention mechanisms. We specifically describe the goals of the system model and algorithm in Section III. The optimization of the edge caching strategy and the two proposed methods are introduced in Section IV. The simulation results and a case study are conducted in Section V. Finally, Section VI summarizes the full work.

### II. RELATED WORK

#### A. Edge Caching

Traditional caching replacement methods have made significant contributions in this field, such as Least-Recently-Used (LRU), Least Frequently Used Dynamic Aging (LFUDA) [8] and Greedy Dual Size Frequency (GDSF) [9]. Many studies sharpened edge caching by combining popular methods with probability. In [10], the authors envisaged proactive caching systems and distributed caching replacement processes based on the popularity of the content. However, this method neglected the reality of user variability and the high diversity of users’ interests to a certain extent.

Furthermore, with the continuous efforts of researchers, artificial intelligent (AI) methods have been used to solve the problem of caching replacement and have played a milestone role. In deep learning (DL), the neural network was used to investigate user inclinations precisely and predict the request mode of BS [11]. Wang et al. paid attention to the design of the distributed caching strategy in [12], considered the cache capacity of the edge server layer, and proposed a distributed caching replacement method according to the popularity of the content to perform collaborative edge caching between BSs, and finally, minimize the transmission cost. Besides, deep reinforcement learning (DRL) was also a methodology that has made some achievements in caching policy optimization. It supported the network objects to learn and build knowledge and interact with the environment. In [13], the authors built an incentive mechanism, which was used for the distributed caching system and optimized via Deep Q-learning (DQN). In some studies, multi-agent learning centred on getting an overall reward [14], which designed a centralized critic network. However, a singular and particular aggregator leads to an additional transmission cost, and adaptability is also one of the challenges.
B. Attention Mechanism

In recent years, the attention mechanism was used for predicting the click-through rate (CTR) in e-commerce [15], [16] by discerning the influential characteristics of other entities to the current commodity or user. Especially in [17], Zhou et al., authors investigated the distribution of influence, which was captured from the historical operation on the present products. Literatures of [18], [19] concentrated on finding the state’s intrinsic composition and relationship in the complex environment of the multi-agents system.

There were many variants of the attention mechanism, such as hard attention [15], but this method did not pay attention to all the input content, but only paid attention to a specific aspect of information. Another variant was the multi-head attention [20], which aimed to allow the model to obtain more levels of information about the sentence from different representation spaces and improved the feature expression ability of the model. It has made contributions in the field of sentiment analysis and machine translation [21], [22]. To this end, we design the edge caching strategy by considering the multi-head attention mechanism for better extracting the underlying system information.

III. System Model

In this section, we introduce the models of the proposed framework, including the system architecture, user request model, and communication model. Then the objective problem is formulated.

A. System Overview

Fig. 2 shows the heterogeneous caching model consisting of various users, BSs and a core network. Let $B = \{ B_1, B_2, \ldots, B_b \}$ denote the BSs, and the users are randomly distributed in the coverage area of the BS $B_b$, which can be donated as $U_b = \{ U_1, U_2, \ldots, U_u \}$, where $U_u = \{ 1, 2, \ldots, u \}$. The union of $C_b$ is a set of cached contents under all BSs denoted as $C = \{ C_1, C_2, \ldots, C_b \}$, and $N = \{ n_1, n_2, \ldots, n_b \}$ to signify the storage capacity of BSs. The backhaul link is the bridge between the core network and each BS. We consider a content pool $F = \{ f_1, f_2, \ldots, f_F \}$ in the core network, which stores all of the contents. Let $req(u_b)$ express the requested content sent by user $u$ to BS $B_b$ via the cellular links. We only consider the cost that the request of adjoining BS is in the last hop (i.e., a two-hop requested scenario). If the request is not available in the neighbor BS $B_b$\(^1\). The request will be sent to the cloud centre through backhaul if the content cannot be satisfied by its neighbor.

As mentioned in the above section, the network is strengthened by arranging BSs in buildings, such as shopping malls, schools, hospitals, etc., or on outdoor facilities (e.g., parks and traffic lights.) In real life, these scenes are close to each other in various combinations, and the distribution of these BSs is also close and intersects. In the proposed caching framework, the communication from a BS to others is composed of the observation value and action of each BS. To describe the problem clearly in the modelling, we only consider the transfer between neighbors of one hop in the part of communication between BSs. Observation $o_b$ and action $a_b$ of BS $B_b$ thereafter are transferred between it and the one-hop neighboring BSs of it.

B. User Request Model

As verified in [23], Mandelbrot-Zipf (MZipf) is usually used to model the distribution of content popularity as a basis of request sending and transmission in the caching process. In our model, the generation of user request content is also derived from MZipf distribution. Motivated by this, the probability that the content $f$ is requested by users in BS $B_b$ is defined by:

$$\omega_{fb} = \frac{(R_f + q)^{-\alpha}}{\sum_{i=1}^{F} (R_i + q)^{-\alpha}},$$  \hspace{1cm} (1)$$

where $R_f$ represents the popularity rank of content $f$ in descending order in local. Among the parameters, $\alpha$ is the skewness factor and $q$ donates the plateau factor.

C. Communication Model

The communication process mainly includes three processes, as shown in Fig. 2. Process ① indicates the communication between the local BS and user which is established by the cellular links. If the local BS $B_i$ stores the requested content $f$, it can provide the quick content access service for user $u$, thus, the cost of this process is represented as $C_{ij}^B$. However, it is impractical to put all the content in the cache of the local BS $B_i$, the request will be sent to another neighbor BS. This process is illustrated by ②, and the transmission cost $C_{ij}^{BS}$ between the current BS $B_i$ and the neighbor BS $B_j$ is measured by the distance [24], expressed as $C_{ij}^{BS} = \nu d(B_i, B_j)$, $i, j \in \{1, 2, \ldots, b \}$, where the non-negative constant $\nu$ is the cost coefficient, and $d(B_i, B_j)$ is the network distance between the two BSs $B_i$ and $B_j$.

If the neighbor BSs cannot provide the service for the required content through the process ③, i.e., the backhaul link, the indicated request will be sent to the core network. According to the experience of existing research [25], we use a positive parameter $\mu$ to represent the unit transmission cost in the backhaul link. Related to formula (1), the request of

\(^1\)We consider the BSs are connected by fibres in the networks.
the user for content $f$ is not available\(^2\) in BS $B_b$ and the probability of the request to the core network is:

$$\omega_{fbh} = 1 - \omega_{fh}.$$  

(2)

The transmission cost of backhaul is:

$$C_{B_f} = \mu_\omega \omega_{fbh}.$$  

(3)

The main costs incurred during the entire content request process include not only transmission costs. It also includes the caching replacement cost, which is closely related to the size of the content [26]. The overhead incurred during the caching replacement is defined as follows:

$$C_{\text{cache}} = \delta (c_r - c_c), r, c \in \{1, 2, ..., N\}.$$  

(4)

where the content replacement factor $\delta > 0$, $c_c$ and $c_r$ respectively represent the size of two contents before and after the replacement.

<table>
<thead>
<tr>
<th>Notation</th>
<th>Meaning</th>
</tr>
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<tbody>
<tr>
<td>$B$</td>
<td>Set of BSs</td>
</tr>
<tr>
<td>$n_i \in N$</td>
<td>Storage capacity of BS $B_i$</td>
</tr>
<tr>
<td>$\mathcal{F}$</td>
<td>Content pool</td>
</tr>
<tr>
<td>$U_b$</td>
<td>User set under the coverage of BS $B_b$</td>
</tr>
<tr>
<td>$C_b$</td>
<td>Cached content set of BS $B_b$</td>
</tr>
<tr>
<td>$\text{req}(u_b)$</td>
<td>Requested content sent by user $u$ to BS $B_b$</td>
</tr>
<tr>
<td>$S_b$</td>
<td>State space of BS $B_b$</td>
</tr>
<tr>
<td>$A_b$</td>
<td>Action space of BS $B_b$</td>
</tr>
<tr>
<td>$a_b$</td>
<td>Current observation state of BS $B_b$</td>
</tr>
<tr>
<td>$a_o$</td>
<td>Current action made by AC agent</td>
</tr>
<tr>
<td>$\omega_{fh}$</td>
<td>Probability that the content $f$ is requested</td>
</tr>
<tr>
<td>$\alpha, q$</td>
<td>Skewness factor, plateau factor</td>
</tr>
<tr>
<td>$C_{\text{hit}}^{B_i}$</td>
<td>Cost when user’s request satisfied by local BS $B_i$</td>
</tr>
<tr>
<td>$C_{\text{hit}}^{BS}$</td>
<td>Cost when fetch content from neighbor BSs</td>
</tr>
<tr>
<td>$C_{\text{hit}}^{BN}$</td>
<td>Cost when fetch content from Core network</td>
</tr>
<tr>
<td>$r_i$</td>
<td>Reward which used for training RL model</td>
</tr>
<tr>
<td>$\theta_i$</td>
<td>Actor’s network parameter</td>
</tr>
<tr>
<td>$d$</td>
<td>Numbers of feature</td>
</tr>
<tr>
<td>$e$</td>
<td>Embedding feature used for critic</td>
</tr>
<tr>
<td>$Q$</td>
<td>Query embedding for different feature</td>
</tr>
<tr>
<td>$k_{ij}$</td>
<td>Key value for calculating attention weight</td>
</tr>
<tr>
<td>$\zeta_{d, j}$</td>
<td>Calculated attention weight</td>
</tr>
<tr>
<td>$v_{ij}$</td>
<td>Value that contain caching information</td>
</tr>
<tr>
<td>$\psi_o$</td>
<td>Parameter of BS $B_b$ target critic network</td>
</tr>
<tr>
<td>$\gamma$</td>
<td>Decay factor of RL model</td>
</tr>
</tbody>
</table>

D. Problem Formulation

In this part, we introduce the goals which are obtained based on the details and limitations of the modeling described in the past section. The objective of the proposed framework is to keep the costs of the system as low as possible. The system saved costs can be obtained as follows.

We use $G$ to represent the revenue of the entire system:

$$G = \sum_{B_i \in \mathcal{B}} \sum_{u \in U_i} (\text{req}(u)C_{B_f}^{u}) + (1 - \text{req}(u)) \Delta).$$  

(5)

The $\Delta$ in the formula is expressed as the following:

$$\Delta = \left\{ \begin{array}{ll}
\sum_{j \neq i} (\text{req}(u)) (C_{B_f}^{B_i} - C_{B_f}^{BS}) f_{\text{cache}}(u)) & \text{if } \sum_{j \neq i} \text{req}(u) \neq 0 \\
0, & \text{otherwise}
\end{array} \right.$$  

(6)

The parameter $\text{req}(u)$ means the content requested by the user $u$, $\text{req}(u)$ represents whether the content requested by $u$ is stored in the BS $B_i$, and $\text{req}(u) = 1$ means that the content $f$ has been cached, and there are no replacement actions required in this case. If $\text{req}(u) = 0$, the modified content needs to be satisfied through other processes, e.g., directly from neighbour BSs or the core network.

Maximizing the reduction of the transmission and caching cost is the goal of our work:

$$\max \sum_{B_i \in \mathcal{B}} \sum_{u \in U_i} (\text{req}(u))C_{B_f}^{u}$$.  

(7)

s.t. $\sum_{u \in U_i} \text{req}(u) \leq n_i$.

(8a)

$$d(B_i, B_j) > 0,$$  

(8b)

$$\nu, \mu, \delta \geq 0.$$  

(8c)

The constraint shown in (8a) indicates that the sum of the size of the cached content is limited by the storage capacity of the BS. Furthermore, (8b) describes the parameter table of the token replacement behavior, (8c) and (8d) ensure the positive value of the distance between BSs, BS-BS cost coefficient, backhaul unit transmission cost, and replacement factor.

IV. ATTENTION-BASED MULTI-AGENT CACHING REPLACEMENT STRATEGY

Each BS has its own service area, in which users have different preferences and relationships with others. Meanwhile, BSs communicate with each other and cooperate. To achieve an efficient caching strategy, we use the influence and information between BSs to a greater extent. The actor-critic algorithm is an effective combination of strategy-based and value-based methods. We propose a multi-agent framework based on actor-critic to solve the edge caching problem.

A. Multi-Agent Actor-Critic

In our architecture, each BS elaborates one actor-network and one critic-network, where different users have their own interests and preferences. Meanwhile, due to the difference in geographic spaces and social functions as well as the heterogeneous services, BSs may have diverse influences on the neighbors. As shown in Fig. 3, for the BS $B_1$ in the blue area, the BS $B_2$ in the dark orange area has more influence on its caching decision compared with $B_3$ in the light orange area. Agents can also learn from other BSs which help make the intelligent caching decision. In this way, we assume that the
caching state can be passed among BSs for better information (e.g., the size/type of cached content) interactions.

**State Space:** We use $S = \{S_1, S_2, ..., S_b\}$ to describe the state space. The state space of each agent is made up of the cached list state and the information of the cached content, for example, the size and the type of the cached content.

**Action Space:** Let $A = \{A_1, A_2, ..., A_b\}$ represent the action space of all BSs, where $A_i$ means the action space of $B_i$. The size of the action space of each BS is equal to the amount of content currently cached plus one as the agent can also choose the action that does not replace existing content in the current caching list.

**Reward:** Based on the definition system model, we design the system reward as:

$$ R_i = C_{B1}^{BS} - C_{B1}^{cache} - C_{B1}^{FC} $$

where the above expression represents the reward of receiving the requested content from adjacent nodes, and the other one is the reward of receiving it from the current BS.

**Observation:** Let the $o = \{o_1, o_2, ..., o_b\}$ denote the observation, agents can obtain the local caching state and the user’s requests under the coverage of its serving area. Besides, as different BSs can establish communication links, agents can also obtain some information about neighboring BSs as a result. In the next subsection, we will introduce the specific temporal and spatial factors involved in the observation.

**Actor Network:** Each local agent observes its serving area’s local state and then makes the caching decision based on its local state. In the actor-network, the optional actions include all currently cached contents and choose action mainly based on its local observation, the policy gradient therein is used to update the parameter. We use $\theta$ to represent the content replacement parameter:

$$ a_i = \pi_{\theta_i}(o_i) $$

where $o_i$ means the local caching states of the BS $B_i$.

**Critic Network:** The decisions of caching actions are selected by the actor-network deployed in each agent based on local information. After the process of actions selection in the actor-network, the critic-network evaluates the expected reward of the policy according to the observation and actions of all BSs through the value function. Furthermore, an attention mechanism is adopted to model the complicated impact of other agents, and we describe our specific method in the following sections.

**B. Critic Network with Attention Mechanism**

We use time and space factors as observations, and mine the influence of BSs based on the critic network with the attention mechanism. As an agent, each BS has its own critic network and can observe the history caching status of its nearby BSs. The whole process can be regarded as a combination of distributed local training and centralized global learning. In order for the critic to learn globally from the information of neighboring BSs, as shown in Fig. 4, we set up three modules of embedding layer, attention layer, and output layer in the critic network.

**Embedding Layer:** The obtained numerical features will be fed into the multilayer perceptron (MLP) in this layer. Because different BSs provide users with a variety of services in different scenarios, resulting in the uncertain number of services provided by each BS. This results in the length of the corresponding encoding vector list is not a fixed value.

As described in the previous subsection, we use one-hot or multi-hot encoding operations to express the types. We know that the one-hot feature is too sparse, and even the encoding of the caching category with multi-hot encoding may be a very sparse vector, which is not suitable for direct input into the subsequent neural network for training. To this end, it is necessary to convert this sparse one-hot vector into a denser embedding vector by connecting it to the embedding layer.

![Fig. 3. Illustration of actor-critic mechanism in each BS.](image-url)

![Fig. 4. Design of the attention-based critic network.](image-url)
The embedding vector $e_b$ of the $i$-th BS is obtained through the embedding operation. Observation of $B_b$ can be embedded and transformed into embedding vector $e_t$, and $e_j$ is defined as the observation of the adjacent BS $B_j$ linked to the current BS via the BS-BS link. Due to the different caching capabilities of BSs and the different sizes of content, the amount of content in each BS may be different, resulting in different lengths of feature vectors. However, a fully connected network can only handle fixed-length input. To solve this problem, the pooling method is generally adopted to map the original features to a fixed-length vector and ensure the immutability of features. In this way, we have the pooling vector $e = pooling(e_1,...,e_j)$ of the neighboring BSs.

**Attention Layer:** The multi-head attention mechanism is deployed in this layer, which uses multiple queries $Q = \{q_1, ..., q_d\}$, where $d$ means the number of kinds of features, to calculate in parallel to select multiple information from the input information, and each attention focuses on a different part of the input information. The different feature subspaces are shown in Table II. The cached content, request content, and BS feature in the table correspond to head\(_{cached}\), head\(_{request}\), and head\(_{BS}\), respectively. The requested feature is similar to the feature of the cached content and will not be repeated. Our critic is not centralized, in each head the query is a kind of character in the current BS. As described in the previous subsection, the head includes the size, type, and request frequency of the cached content, caching capacity, and the users who are under the service of the local BSs. Key $k_{d,j}$ and value $v_{d,j}$ are the information of the adjacent BS $B_j$ on the feature $d$. After completing the process of information input, we get the attention to aspect $d$ of adjacent node $j$, which is named the influence unit $ζ_{d,j}$:

$$ζ_{d,j} = \sum_{j \neq i} \frac{q_{d,j}^T k_{d,j}}{\exp(q_{d,j}^T k_{d,k})},$$

(11)

where $d$ is used to mark the features observed by the current head.

After getting the influence unit $ζ_{d,j}$, we calculate cooperation unit $Att_{d,b}$ through the weighted sum of attention weight $ζ_{d,j}$ and value $v_{d,j}$. In this way, we can obtain the $k_{d,j}$ and $v_{d,j}$ through different linear embedding layers where $k_{d,j}$ is used to calculate weights and $v_{d,j}$ is used to capture features. $Att_{d,b}$ is defined as:

$$Att_{d,b} = \sum_{j \neq b} ζ_{d,j} v_{d,j} .$$

(12)

**Output Layer:** Based on the attention operation of the previous layer, at this stage, we concatenate these weights in columns, and then perform a linear transformation with a new weight matrix $W^a$ to get the final attention output. The concatenate function is used to aggregate the attention of multi-heads:

$$Att_b = concat(Att_{1,b},...,Att_{d,b}) .$$

(13)

The critic network of each BS estimates the action-value function by attention mechanism which can be expressed by:

$$Q_ψ_b(o,a) = σ_b(e_b, Att_b) ,$$

(14)

where $ψ$ is the parameter of the critic network, $σ_b$ is an MLP layer, and the $ψ_b$ is the parameter of the target critic network which is associated with the agent $B_b$.

### C. The Algorithm process

Reinforcement learning (RL) is an iterative process. Each iteration needs to be given a strategy evaluation function and update the strategy according to the value function. There are few studies on the convergence of actor-critic, the main reason is that the update has a strong network and status relevance. Therefore, in addition to the original ac network, a target ac network is also established. The target network has the same structure and initialization as the original network. While training parameters of the networks, the target network estimates future actions.

In the RL scheme, the objective is to find the optimal policy. To this end, the Q-function is used to represent the expectation of the total reward that the agent can obtain in the future after taking action $a$ in state $s$:

$$Q_i(o,a) = E_{π_i}\left[\sum_{t=0}^{+∞} γ^t r_{t+1}|o_0 = o, a_0 = a\right] .$$

(15)

To get the optimal $Q$-value, the input corresponding state and all possible actions are required to be traversed and then find the largest $Q$-value. We use equation (14) to approximate equation (15).

The attentive critic network calculates the advantage function to get the advantage of the current action compared to other actions [7], and the advantage function is applied to the gradient update process of individual strategy as follows:

$$A_ψ^i(o,a) = Q_ψ_i(o,a) - \sum_{a_i' \in A_i} π_θ_i(a_i'|a_i)Q_ψ_i(o_i(a_i,a_i'), a_i) .$$

(16)

This additional item is a baseline that considers an additional benefit in a multi-environment interaction, and it does not change the decision-making process of the agent. The calculation of the updated gradient of our policy is shown as:

$$∇_θ_i J(π_θ_i) = E_{π_θ_i}[∇_θ_i log(π_θ_i(\theta_i | a_i))A_i(o,a,ψ_i)] .$$

(17)
The update of the critic network is based on the following loss function:

$$L(\psi) = \mathbb{E}[(y_i - Q_\psi(o, a))^2].$$  \hfill (18)

In the above equation, the $y_i$ is defined as:

$$y_i = r_i + \gamma \mathbb{E}_{\pi_{\theta_i}}[Q^*_i(o', a') - \rho \log(\pi_{\theta_i}(a'_i|o'_i))],$$  \hfill (19)

where $\gamma$ is the decay factor, $\psi_i$ and $\overline{\psi}_i$ are the past parameters before update, $\rho$ determines the balance between maximizing entry and rewards [27].

The parameter $\overline{\psi}$ of critic network and $\theta$ of actor-network are updated after each iteration:

$$\overline{\psi} = \psi - lr_c \nabla L(\psi_i),$$ \hfill (20)

$$\theta = \theta - lr_a \nabla J(\pi_{\theta_i}),$$ \hfill (21)

where $lr$ is the learning rate, which is set in the experiment according to the actual situation.

Note that the above two equations are used to increase the randomness of our strategy which can decentralize the probability of each action, in case the agent just chooses one action. The whole process is as shown in Algorithm 1. Like in [4], the time complexity of back-propagation is $O(abN^2)$, in which $a$ and $b$ are the number of layers and unit, respectively.

### Algorithm 1 The procedure for caching by NAC

1: Initialization:
   Parameter of each actor network $\theta_i$
   Parameter of each critic network $\psi_i$
   Reset the parallel environments with $B_h$ agents
2: $T_{update} \leftarrow 0$
3: for $t = 1, T$ do
4:   each BS receives requests $R_t = \{r_1, r_2, \ldots, r_b\}$
5:   obtain BSs’ observations $o_t = \{o_1, o_2, \ldots, o_b\}$
6:   each BS selects by the policy $\pi_{\theta_i}(a_b|o_b)$
7:   each BS sends action and state to adjacent BSs
8:   $T_{update} = T_{update} + 1$
9:   if $T_{update} \geq$ min steps per update then
10:      calculate $Q_{\psi_i}(a_{1..b}, a_{1..b})$ for all BSs
11:      calculate $a'_i = \pi_{\overline{\psi}_i}(o'_i)$ using target policies
12:      calculate $Q_{\psi_i}(a'_i, a'_{1..b})$ for all BSs
13:      update critic using $\nabla L_Q(\psi_i)$
14:      calculate $a_i = \pi_{\overline{\psi}_i}(o'_i)$ for all BSs
15:      calculate $A^0_{\theta_i}(o, a)$ for all BSs
16:      update policies using $\nabla \theta_i J(\pi_{\theta_i})$
17:      update the parameters for all BSs
18:      update state of each BS
19:   $T_{update} \leftarrow 0$
20: end if
21: end for

### D. An Information-based Engineering Solution

Considering the deployment cost of neural networks embedded in all scenarios, inspired by the core idea of the attention mechanism, we abstract the information of the Neighbors’ Influence (NI) and design a heuristic method by improving the traditional LFU, named NI-LFU, from the engineering perspective.

Based on the temporal locality of traffic, each content that has been cached is assigned a value $F(f, b) = \sum_{j=1}^{k} m^j, j \in \{1, 2, \ldots, k\}$, where $m^j$ means if the current BS meets the requested content $f$, and $k$ is the number of requests since the content $f$ was cached. Taking into account the impact of multi-granularity unit time and surrounding neighbor nodes, we change the value as:

$$F_{NI}(f, b) = \frac{1}{T_{update}} \sum_{i=0}^{B_h} (F_m(f, i) + F_w(f, i) + F_d(f, i)),$$

where $F_m(f, i), F_w(f, i), F_d(f, i))$ indicates the request frequency in the different three modes under the service range of the BS $B_i$ respectively. The long-term frequency of request content $f$ is defined as its requested times in a month $F_m(f)$, the mid-term frequency is its requested times in a week $F_w(f)$, and the short-term frequency is represented by the content request frequency in a day $F_d(f)$. Let $F_{NI}(f, b)$ denote the elimination priority on which content will be replaced. The workflow of the NI-LFU algorithm is shown in Fig.5.

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**Fig. 5. The workflow of NI-LFU algorithm.**

### V. Case Study and Experiment

In this section, we design a small scratch for the adaptability of heterogeneous environments and evaluate the performance of the proposed framework based on the simulation results.

#### A. Simulation Configuration

In order to reflect the feasibility of the theory, we designed a small-scale simulation experiment, and use the stochastic
geometry \([28]\) to model the potential correlation between BSs and the service area of BSs in a multi-agents environment. Stochastic geometry has proven to be a powerful tool for evaluating wireless network performance. Fig. 6 simulates the relationship between BSs, where the red triangles indicate BSs and the blue stars are users.

The total number of contents is \(F = 10000\), and different contents have a normalized size of different values which range between 1MB and 10MB. In order to conduct the case study experiments, we set the total number of BSs as 4, and the cache size of each BS is set to 100MB as default. Users are distributed in the serving range of each BS randomly. In addition, we use a parameterized network to solve the problem, and we confirm some of the parameters (i.e., the learning rate of actor-network \(l_{r_a}\), the learning rate of critic \(l_{r_c}\)) involved in the experiment.

We compare our methods with the following state-of-the-art baseline algorithms:

1) **Actor-critic (AC):** A TD method that has a separate memory structure to explicitly represent the policy independent of the value function \([29]\).

2) **Deep Q Networks (DQN):** In addition to the difference in network structure, the method differs from the proposed framework in that it is based on local characteristics and does not add influencing factors in the environment of neighbor nodes. DQN and AC are mainly used to compare with NAC to verify the stability of the NAC algorithm \([30]\).

3) **First-Input-First-Output (FIFO):** FIFO is one of the oldest caching methods. This algorithm always deletes the content with the longest time \([5]\).

4) **Least-Frequently-Used (LFU):** In the LFU algorithm, the least frequently used content is knocked out \([31]\).

5) **Least-Recently-Used (LRU):** LRU is a common caching replacement algorithm that eliminates the most recently unused content \([5]\).

6) **Least-Frequently-Used-Dynamic-Aging (LFUDA):** LFUDA builds upon simple LFU by accommodating shifts in the set of popular objects in the cache \([31]\).

7) **Greedy-Dual-Size-Frequency (GDSF):** A strategy based on the value relationship of all cache objects in the cache node \([9]\).

8) **Neighbor-Influence-Least-Frequently-Used (NILFU):** From the perspective of information, NI-LFU is an engineering solution formed by the idea of our approach. It combines the request frequency in the current area and the neighbor environment via multiple angles to make replacement decisions.

We take the following indicators to evaluate the performance of different methods:

**Hit Rate:** it is the ratio of hits and total requests. Hit means that the user can get the required data directly through the cache. At the same time, the performance index we use when conducting parameter selection experiments is also a hit rate.

**Average Transmission Delay:** it is calculated based on the delay between the cloud data centre and the BS, between BSs, and between the user and BS.

**Traffic Offload Ratio:** it means the rate of BS offloaded traffic to the total traffic.

### B. Case Study

A key challenge in highly heterogeneous services is making full use of complex spatial and temporal characteristics \([32]\). We designed this caching replacement optimization strategy with the attention mechanism and used artificial intelligence technology to perceive the environment of neighbor BSs.

To prove the practicability of the proposed framework in a heterogeneous environment, we design a small demo to compare the algorithm’s performance in a heterogeneous environment and a homogeneous environment.

![Fig. 7. Distributions of Content popularity for 4 BSs.](image)

In this case, as expressed in Fig. 6, we consider 4 BSs, and the storage size of BSs is [300, 340, 380, 420]. The user number of each BS is [10, 7, 5, 8], and the initial state of the user is randomly generated. At the same time as agents, their training parameters are also different, the experience pool sizes of these agents are [2000, 2500, 1500, 2000], the episode numbers of BSs are set as [200, 150, 250, 150], the batch sizes are [128, 128, 256, 256].

We have the following two cases, i.e., heterogeneous and homogeneous cases. (1) For the heterogeneous case, we simulate heterogeneity of different services by setting the MZipf factor and plateau factors in equation (1), the popularity distributions are shown as dotted lines in Fig. 7. (2) For the homogeneous case, the content popularity follows the same MZipf, shown...
as the solid line in Fig.7, we set the value of $q$ and $\alpha$ based according to the reference [33] which originally proposed the Mandelbrot–Zipf model for content popularity.

In the homogeneous environment, the agent can easily obtain the network information while the spatial and temporal characteristics in heterogeneous are dynamic and more complex. As shown in Fig. 8(a) and Fig. 8(c), the hit rate and traffic offload ratio of the two kinds of environments is relatively close. It indicates that the proposed NAC can achieve similar performance in terms of hit rate and traffic offload ratio. This is mainly because the proposed framework learns the information from the neighbor BSs which improves the network performance. Besides, the curves interact several times due to the dynamic environment changes. Especially in Fig. 8(b), when the cache size is 180, the average delay reduces by 2.1% at best in the heterogeneous situation compared with the homogeneous one. This shows our method has a great capacity for neighborhood awareness and adaptability in heterogeneous scenarios.

C. Parameters Settings

![Fig. 9. Performance of the hit rate under different learning rates of actor-network.](image)

Before the performance comparison simulation, we first compare and select the experimental parameters of our method. We use the training round as the abscissa to show the effect of the different parameters, Fig.9 compares the performance with different learning rates $lr_a$ of the actor-network. The comparative learning rate is $lr_a = [0.01, 0.05, 0.1]$. To maintain the algorithm’s stability and efficiency, we defined the learning rate $lr_a = 0.05$ as an empirical value.

We demonstrate the network performance of the hit rate in the following Fig.10. We compare the performance for hit rate under different learning rates of critic network $lr_c = [0.01, 0.05, 0.1]$. We can see that if $lr_c = 0.05$, the proposed framework produces the best hitting effect. Therefore, $lr_c = 0.05$ is used as an algorithm parameter of the NAC in our simulation.

![Fig. 10. Performance of the hit rate under different learning rates of critic network.](image)

D. Comparison of Different Training Methods

We compared NAC, AC, DQN, FIFO, LFU, LRU, LFUDA, GDSF, and NILFU under different BSs’ buffer sizes. The corresponding simulation results are shown in Fig. 11-Fig. 13. Fig. 11 shows the performance of the hit rate. The hit rate of NACE is obviously the highest. Compared with several commonly used traditional methods, it has increased by 26.81% (FIFO), 26.14% (LFU), 29.04% (LRU), 22.62% (LFUDA), 20.45% (GDSF), 11.23% (AC) and 27.43% (DQN) under different parameters. Moreover, although the NILFU we designed is not based on machine learning, after adding the influence of neighboring cells, compared with other traditional
non-learning caching replacement algorithms, the advantage
hit rate of it is also obvious. Compared with AC, DQN, FIFO,
LFU, LRU, LFUDA, and GDSF, the hit rate increased by
1.26%, 16.00%, 15.44%, 14.83%, 17.47%, 11.31% and 9.65%,
respectively.
As shown in Fig. 12, compared with several traditional
replacement methods, NAC can effectively reduce the delay
by 0.030s (AC), 0.086s (FIFO), 0.190s (LFU), 0.086s (LRU),
0.072s (LFUDA), 0.072s (GDSF). Although NAC’s perfor-
mance in latency is slightly inferior to DQN, it is better in
terms of hit rate and traffic unloading. And from the figure,
we can see that compared to DQN, as the cache capacity increases,
the performance of NAC will improve. As an algorithm based
on a value function, DQN itself is not easy to find the greatest
Q-value among many values when facing a larger action space.
With the enhancement of the caching capacity, the action space
becomes larger. This shortcoming is also exposed, and on the
other hand, it also illustrates the advantages of NAC. And the
NILFU average delay performance is excellent. As shown in
Fig. 12, the average delay of NILFU is reduced by 0.050s,
0.154s, 0.053s, 0.037s, and 0.037s compared to FIFO, LFU,
LRU, LFUDA, and GDSF, evenly and respectively. When the
BS cache capacity is large to some degree, NILFU can be
equal to or even surpass the AC algorithm. This also fully
proves that NILFU not only uses neighborhood information
to improve accuracy but also can overcome some of the
disadvantages of complex algorithms.

At the same time, as depicted in Fig. 13, NAC and NILFU
also show good advantages in offload. From the figures, we can
see that with the enhancement of the cache capacity, the tra-
tditional AC does not have a higher performance improvement
than NAC. With the enhancement of caching capabilities, the
content cached by BSs has a higher diversity, only considering
local information, so AC cannot make a good global decision.
Agents do not consider the dependency of the surrounding
node cache, so its performance is relatively poor. It can be
seen from the resulting graph that the broken line of the
neighborhood-aware method will be more smooth and more
stable. But it will have relatively large fluctuations in the
traditional AC method. This shows to a certain extent that
steady.

VI. CONCLUSION
In this paper, we have proposed an edge caching framework
called NAC by considering the neighborhood-aware informa-
tion, where the Actor-Critic method is employed to solve
the complicated edge caching problem. The AC agent can
make efficient caching decisions after gathering all the global
and local information, leading to excellent backhaul traffic
reduction and QoS/QoE improvement. Each AC agent contains
two networks: the actor-network chooses actions according
to the local information. The critic network evaluates the
impact of surrounding BSs. What is more, to integrate the
information transmitted between BSs, we have employed a
multi-head attention mechanism to quantify the impact of
neighboring BSs on the current node with each head capturing
the characteristics under a subspace. Besides, we also propose
a non-learning caching algorithm NILFU, which can achieve
considerable performance in some scenarios without the use of
AI services. In the section of experiments, the proposed NAC
algorithm shows the feasibility and effectiveness compared to
several existing caching algorithms. In the future, We will continue to improve our algorithm and architecture based on the actual situation.

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technical papers in IEEE JSAC, TCC, ToN, TWC, IoTJ, COMST, TMM, INFOCOM, ICDCS and so on. He has received the best paper awards of IEEE ICC, ICPADS, and in 2017, he was the recipient of the “IEEE ComSoc Fred W. Ellersick Prize”, and in 2022, he received the “IEEE ComSoc Asia-Pacific Outstanding Paper Award”.

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